

# CS 2336

# Discrete Mathematics

## Lecture 11

Sets, Functions, and Relations: Part III

# Outline

- What is a Relation ?
- Types of Binary Relations
- Representing Binary Relations
- Closures

# Cartesian Product

- Let A and B be two sets

The **cartesian product** of A and B, denoted by  $A \times B$ , is the set of all ordered pairs

$$\{ (a, b) \mid a \in A \text{ and } b \in B \}$$

- $A = \{ 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, J, Q, K \} =$  all ranks

$$B = \{ \spadesuit, \heartsuit, \diamondsuit, \clubsuit \} = \text{all suits}$$

→  $A \times B =$  all 52 cards in a deck

$$= \{ (1, \spadesuit), (2, \spadesuit), (3, \spadesuit), \dots, (J, \clubsuit), (Q, \clubsuit), (K, \clubsuit) \}$$

# Cartesian Product

- Let  $A_1, A_2, \dots, A_k$  be  $k$  sets

The **cartesian product** of  $A_1, A_2, \dots, A_k$ , denoted by  $A_1 \times A_2 \times \dots \times A_k$ , is the set of all ordered pairs  $\{ (a_1, a_2, \dots, a_k) \mid a_j \in A_j \text{ for all } j = 1, 2, \dots, k \}$

- Let  $A_j =$  the set  $\mathfrak{R}$  of real numbers, for all  $j$   
→  $A_1 \times A_2 \times A_3 =$  the 3-d Euclidean space  $\mathfrak{R}^3$

# Relation

A **binary relation from A to B** is a subset of the cartesian product  $A \times B$

- Example:

$A = \{ 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, J, Q, K \}$

$B = \{ \spadesuit, \heartsuit, \diamondsuit, \clubsuit \}$

Spades =  $\{ (1, \spadesuit), (2, \spadesuit), (3, \spadesuit), \dots, (Q, \spadesuit), (K, \spadesuit) \}$

➔ Spades is a binary relation

# Relation

A **k-ary relation** is a subset of a cartesian of k sets

- Example:

$$L = \{ (x, y, z) \mid 2x + 3y + z = 0, \text{ and } x, y, z \in \mathfrak{R} \}$$

➔ the line L is a ternary relation of the space  $\mathfrak{R}^3$

# Binary Relation

- In the remaining of this lecture, we focus on a special type of relations :

the binary relation from a set  $A$  to  $A$

- Such a relation is called a binary relation **on**  $A$
- Example :  $A =$  the set of integers

$$R = \{ (a, b) \mid a - b \geq 10 \}$$

# Types of Binary Relations

A binary relation  $R$  on  $A$  is said to be **reflexive** if  $(a, a) \in R$  for every  $a \in A$

- Which of the following relations are reflexive ?
  - $R = \{ (a, b) \mid a - b \geq 10, a, b \text{ are integers} \}$
  - $S = \{ (a, b) \mid a \leq b, a, b \text{ are integers} \}$
  - $T = \{ (a, b) \mid a < b, a, b \text{ are integers} \}$
  - $U = \{ (x, y) \mid x \text{ and } y \text{ are on the same weekday, } x, y \text{ are days in April 2013} \}$

# Types of Binary Relations

A binary relation  $R$  on  $A$  is said to be **symmetric** if  $(a, b) \in R$  implies  $(b, a) \in R$

- Which of the following relations are symmetric ?
  - $R = \{ (a, b) \mid a - b \geq 10, a, b \text{ are integers} \}$
  - $S = \{ (a, b) \mid a \leq b, a, b \text{ are integers} \}$
  - $T = \{ (a, b) \mid a < b, a, b \text{ are integers} \}$
  - $U = \{ (x, y) \mid x \text{ and } y \text{ are on the same weekday, } x, y \text{ are days in April 2013} \}$

# Types of Binary Relations

A binary relation  $R$  on  $A$  is said to be **antisymmetric** if  $(a, b) \in R$  implies  $(b, a) \notin R$  unless  $a = b$

- Which of the following are antisymmetric ?
  - $R = \{ (a, b) \mid a - b \geq 10, a, b \text{ are integers} \}$
  - $S = \{ (a, b) \mid a \leq b, a, b \text{ are integers} \}$
  - $T = \{ (a, b) \mid a < b, a, b \text{ are integers} \}$
  - $U = \{ (x, y) \mid x \text{ and } y \text{ are on the same weekday, } x, y \text{ are days in April 2013} \}$

# Types of Binary Relations

A binary relation  $R$  on  $A$  is said to be **transitive** if  $(a, b), (b, c) \in R$  implies  $(a, c) \in R$

- Which of the following are transitive ?
  - $R = \{ (a, b) \mid a - b \geq 10, a, b \text{ are integers} \}$
  - $S = \{ (a, b) \mid a \leq b, a, b \text{ are integers} \}$
  - $T = \{ (a, b) \mid a < b, a, b \text{ are integers} \}$
  - $U = \{ (x, y) \mid x \text{ and } y \text{ are on the same weekday, } x, y \text{ are days in April 2013} \}$

# Representing Binary Relations

- Let  $A$  be a finite set
- The binary relation on  $A$  can be conveniently represented in two different ways :

- Method 1: Matrix Form

$$A = \{ 1, 2, 3, 4 \}$$

$$R = \{ (1,1), (1,2), (2,3), (2,4), (3,4), (4,2) \}$$

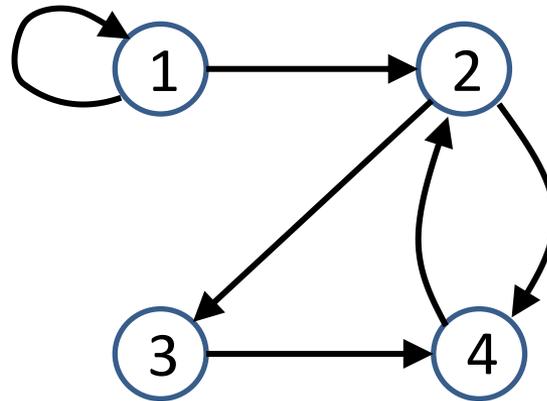
	1	2	3	4
1	✓	✓		
2			✓	✓
3				✓
4		✓		

# Representing Binary Relations

- Method 2: Directed Graph

- $A = \{ 1, 2, 3, 4 \}$

$$R = \{ (1,1), (1,2), (2,3), (2,4), (3,4), (4,2) \}$$



# Closures

- Given a binary relation  $R$ , we may obtain a new relation  $R'$  by adding items into  $R$ , such that  $R'$  will have certain property

- Example :

$$R = \{ (1,1), (1,2), (2,3), (2,4), (3,4), (4,2) \}$$

If we add  $(2,2)$ ,  $(3,3)$ , and  $(4,4)$  into  $R$ , the resulting relation will be reflexive

# Closures

- Let  $R$  be a binary relation

The smallest possible relation  $R'$  that contains  $R$  as a subset, such that  $R'$  has a **property  $P$** , is the **closure of  $R$  with respect to  $P$**

	1	2	3	4
1	✓	✓		
2			✓	✓
3				✓
4		✓		

$R$

	1	2	3	4
1	✓	✓		
2		✓	✓	✓
3			✓	✓
4		✓		✓

Reflexive closure of  $R$

# Closures

- What is the transitive closure of  $R = \{ (1,1), (1,2), (2,3), (2,4), (3,4), (4,2) \}$  ?

	1	2	3	4
1	✓	✓		
2			✓	✓
3				✓
4		✓		

R

	1	2	3	4
1	✓	✓	✓	✓
2		✓	✓	✓
3		✓	✓	✓
4		✓	✓	✓

Transitive closure of R

# Finding Transitive Closure

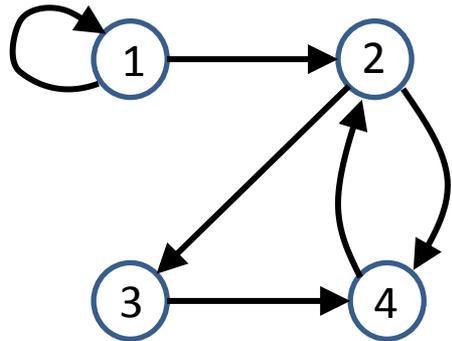
- Getting the transitive closure seems difficult  
Is there a systematic way to get this ?
- Consider the directed graph representation  
→  $R =$  all pairs of vertices where one can reach the other in 1 step
- We can (not easily) show that for vertices  $x$  and  $y$ ,

$x$  can reach  $y$  in the graph  
 $\Leftrightarrow (x, y)$  is in the transitive closure of  $R$

# Finding Transitive Closure

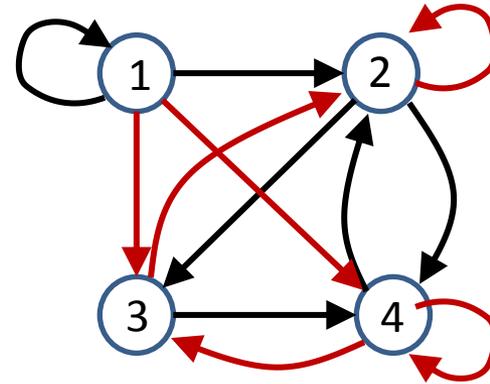
- Let  $R^k =$  all pairs of vertices where one can reach the other in exactly  $k$  steps
  - ➔  $R = R^1$
- We can repeatedly obtain  $R^2, R^3,$  and so on, until we cannot add any new edges
  - ➔ the resulting graph corresponds to the transitive closure of  $R$

# Finding Transitive Closure

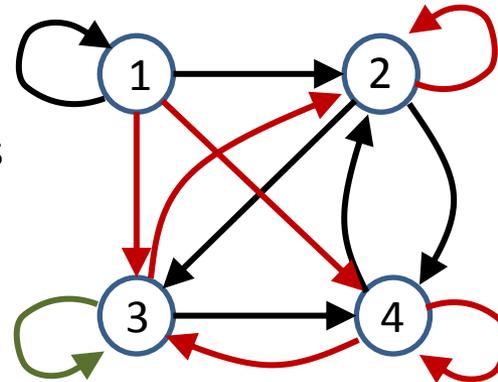


R

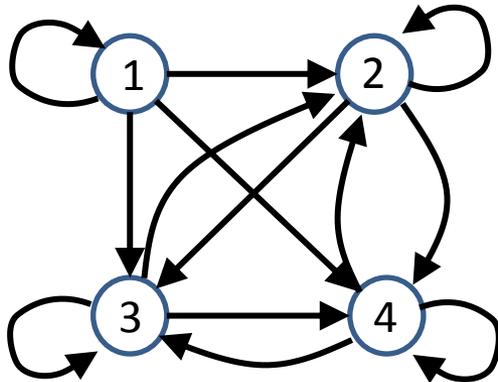
adding  $R^2$



adding  $R^3$



no new edges



Transitive closure of R

# Finding Transitive Closure

- Apart from the previous method, there are faster ways to compute the transitive closure
- If the matrix form is given, and  $|A|=n$ 
  1. Recursive doubling algorithm :  $O(n^3 \log n)$  time
  2. Floyd-Warshall algorithm :  $O(n^3)$  time
    - Code is super simple : Only 3 for loops !